

Integrating Algebraic Dynamic Programming in Combinatorial Optimization

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December 2nd, 2016



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Dynamic Programming & Metaheuristics

Hybrid Metaheuristics often depend on Dynamic Programming for ...

- ... solving **subproblems** e.g. packing, shortest path
- ... enhancing **neighbourhood search** e.g. Dynasearch
- ... improving **recombination** operators in GAs e.g. memetic algorithms
- ... **decoding** solutions e.g. permutation encodings

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Is something wrong with Dynamic Programming?

Algebraic Dynamic Programming (ADP)

Alternative view on Dynamic Programming (Giegerich et al., 2002)

- Formal grammar defines the search space by **decomposition**
- Separates evaluation from search space declaration
- Works for sequence data (strings)—originally intended for bioinformatics
- Extension for set/general data structures available (Siederdisen et al., 2014/15)

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Whistle: a new solver framework for ADP

- targeted for general combinatorial problems
- intended for integration in heuristics

Parts of an Algebraic Dynamic Program

- Set of **indexed terminal symbols**
 - Represents atomic objects of a solution
- Set of **indexed non-terminal symbols**
 - Each non-terminal is a **DP table**
 - addressed by the indices
 - Each indexed non-terminal represents a **state/compound object**
- Set of **production/decomposition rules**
 - Describes the **search space**
 - **Quantifiable**
 - Different types of **constraints**

Motivating Example: Knapsack

Given set of items $i \in \mathcal{I}$ and knapsack of max. weight Q

- S ... Optimally packed knapsack
- $B_{i,q}$... Knapsack of weight at most q with item i considered last
 - i ... integer
 - q ... real-valued
- π_i ... Item i

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Decomposition Grammar

$$\begin{aligned} S &\rightarrow \pi_i B_{i, Q-w_i} \geq 0 && \forall i \in \mathcal{I} \\ B_{i,q} &\rightarrow \pi_j B_{j, q-w_j} \geq 0 && \forall j \in \mathcal{I}, [i < j] \\ &| \epsilon \end{aligned}$$

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Evaluation Algebra σ_{value}

$$\begin{aligned} \sigma_{\text{value}}(S) &= \text{value}[\#1] + \sigma_{\text{value}}(\#2) \\ \sigma_{\text{value}}(B_{i,q}) &= \text{value}[\#1] + \sigma_{\text{value}}(\#2) \\ &| 0 \end{aligned}$$

Dominance: $A \prec B \equiv \sigma_{\text{value}}(A) < \sigma_{\text{value}}(B)$

Heuristic Extensions

Search engines

Original ADP approach uses a fixed search order ...

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- ... for proven-optimality nothing else is needed

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Flexible search orders ...

- ... separate search from search space declaration
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- ... have complexity benefits for some problems

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Whistle supports different **search engines**

- Depth-First Search
- Greedy Search
- A*
- ...

Index propagation

Original ADP

- Not explicit indices (by default)
- Automatic deduction
- Restricted to sequence/set data
- No index errors

Whistle ADP

- Explicit indices (by default)
- No automatic deduction
- Index propagators:
 - Sequence data
 - Cyclic permutations
 - Resource usage
 - ...
- Less index errors
- More flexibility

Partial Invalidation

DP approaches can be embedded in heuristics ...

Improvement heuristics ...

- ...change parts of a solution
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Partial Invalidation ...

- ... keeps track of dependencies of table cells
- ... allows for invalidation of parts of a table
→ on basis of changed terminal symbols
- ... can reuse remaining information

Shadowing

In **Genetic Algorithms** solution candidates ...

- ... depend on their parents
- ... can reuse their information

Shadowing of table cells allows to redirect table access
→ less recomputation

Examples

Shortest Path

Given a graph $G = (V, A)$

- $S_{s,t}$... Shortest path from s to t
- $P_{s,X,t}$... Path from s to t with unvisited nodes X
 - s, t ... integer
 - X ... set
- $a_{i,j}$... Arc from i to j

$$\begin{aligned} S_{s,t} &\rightarrow P_{s, V \setminus \{s,t\}, t} \\ P_{s,X,t} &\rightarrow \begin{array}{l} a_{s,x} P_{x, X-x, t} \\ | \\ a_{s,t} \end{array} \quad \forall x \in X, \left[\begin{array}{l} (s, x) \in A \\ (s, t) \in A \end{array} \right] \end{aligned}$$

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Shortest path is not expressible without set semantics!

Shortest Path with Resource Constraints

Given graph $G = (V, A)$ and k resource capacities $Q^{(k)}$

$$S_{s,t} \rightarrow P_{s, V \setminus \{s,t\}, t, Q^{(k)}}$$

$$P_{s, X, t, q^{(k)}} \rightarrow a_{s,x} P_{x, X - x, t, q^{(k)} - r_{s,x}^{(k)}} \quad \forall x \in X, [(s, x) \in A] [\forall k : q^{(k)} - r_{s,x}^{(k)} \geq 0]$$

$$\quad | \quad a_{s,t} \quad [(s, t) \in A] [\forall k : q^{(k)} - r_{s,t}^{(k)} \geq 0]$$

Traveling Salesman Problem

Given a graph $G = (V, A)$ visit all vertices in V exactly once

Formalization of the Bellman-Held-Karp algorithm

$$\begin{aligned} S &\rightarrow a_{1,i}P_{i, V \setminus \{1, i, j\}, j} a_{j,1} && \forall i, j \in V, [1 \neq i \neq j] [(1, i) \in A] [(j, 1) \in A] \\ P_{i, X, j} &\rightarrow a_{i,x}P_{x, X-x, j} && \forall x \in X, [(i, x) \in A] \\ &| a_{i,j} && [X = \emptyset] [(i, j) \in A] \end{aligned}$$

New Theoretical Insights

Considering the similarity of Shortest Path and TSP models ...

Why is one significantly harder than the other?

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In some cases ...

- ... table indices can be **relaxed** ... **not symbol indices!**
→ multiple indexed symbols map to the same table cell
- ... indices can be stored in an **amalgamated form**
- symbols with a higher **degree of freedom** are computed
→ then update amalgamated index

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Preconditions are already formalized

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Two possibilities for **Shortest Path** ...

- **Amalgamated set index:** less visited nodes \implies higher degree of freedom
- **Completely relaxed set index:** requires heuristic search order \rightarrow Dijkstra's algorithm

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Not applicable for the **TSP!**

Conclusion

Whistle—ADP for combinatorial optimization

...no need to implement all this by yourself!

- Tailored for **combinatorial optimization** in general
- Written in **Rust** as compiler plugin—C ABI compatible
- Supports **integer, float, and set indices**
- Uses a new **compatibility and dominance** mechanism instead of objective functions
- Supports **Index Propagators** for advanced index deduction:
Sequence data, Cyclic permutations, Resources, ...
- Different **evaluation algorithms**:
Top-down, Bottom-up, Bidirectional (new)
- Supports **different search engines**:
DFS (current), Greedy, A*, Beam-Search
- Supports **Partial Invalidation** and **Shadowing**

Thank you for your attention!