

# Perfect Sequence Covering Arrays and related concepts

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## Notation

General assumption: nonneg. integers  $n \geq k \geq 3$ .

$$S_{n,k} := \{(j_1, \dots, j_k) \in \{1, \dots, n\}^k : i \neq \ell \Rightarrow j_i \neq j_\ell\}.$$

Under consideration: Families  $\mathcal{F} = \{\pi_1, \dots, \pi_d\} \subseteq S_n$  ( $i \neq \ell \not\Rightarrow \pi_i \neq \pi_\ell$ ).

Definition (Kuhn et al. 2012; Spencer 1971: "completely  $k$ -scrambling family")

$\mathcal{A} \subseteq S_n$  is called *Sequence Covering Array (SCA)* of strength  $k$ , if

$$\forall x = (x_1, \dots, x_k) \in S_{n,k} \exists \pi \in \mathcal{A} \quad \pi^{-1}(x_1) < \dots < \pi^{-1}(x_k). \quad (*)$$

$\pi$  in (\*) covers  $x$ .

E.g.  $n = 4, k = 3$ :

$$A = \begin{array}{cccc} 3 & 1 & 2 & 4 \\ 1 & 3 & 4 & 2 \\ 2 & 3 & 4 & 1 \\ 4 & 1 & 2 & 3 \\ 2 & 1 & 4 & 3 \\ 4 & 3 & 2 & 1 \end{array}$$

$$(x_1, x_2, x_3) = (1, 4, 3)$$

- $N^*(n, k)$  defined as min.  $d \in \mathbb{N}$  s.t.  $\exists$  strength- $k$  SCA over  $n$  symbols consisting of  $d$  permutations
- Bounds for  $N^*(n, k)$  employed for
  - Order theory, combinatorial geometry, graph theory, etc.
  - Combinatorial Software Testing with SCAs (Kuhn et al. 2012)
  - Min-wise Hashing: Rapid document similarity estimation in huge collections (Google, AltaVista)

## Example (SCA with $n = 7, k = 3$ )

EIUSB: Establish internet connection via USB port;  
 CWLAN: Connect to a WLAN;  
 CWCAM: Connect peripheral device webcam;  
 ENBLU: Enable Bluetooth connectivity;

CMOUS: Connect peripheral device mouse;  
 LOGOI: Logout and successively login;  
 SUSWU: Suspend and wakeup

Performed test							Error?
EIUSB	CWLAN	CWCAM	ENBLU	CMOUS	LOGOI	SUSWU	/
ENBLU	CWCAM	CWLAN	EIUSB	SUSWU	LOGOI	CMOUS	/
EIUSB	CMOUS	LOGOI	SUSWU	CWLAN	CWCAM	ENBLU	/
SUSWU	LOGOI	CMOUS	EIUSB	ENBLU	CWCAM	CWLAN	/
CWLAN	CMOUS	EIUSB	CWCAM	ENBLU	LOGOI	SUSWU	/
CMOUS	CWLAN	SUSWU	LOGOI	ENBLU	CWCAM	EIUSB	Yes
CWCAM	LOGOI	EIUSB	CWLAN	ENBLU	CMOUS	SUSWU	/
LOGOI	CWCAM	SUSWU	CMOUS	ENBLU	CWLAN	EIUSB	Yes
ENBLU	SUSWU	EIUSB	CWLAN	CWCAM	CMOUS	LOGOI	/
SUSWU	ENBLU	LOGOI	CMOUS	CWCAM	CWLAN	EIUSB	Yes

## Theorem (Spencer–Hajnal; Radhakrishnan–Spencer)

$$N^*(n, k) = \Theta(\log_2 n) \quad (k \geq 3, \text{ fixed}).$$

- $N^*(n, k) \geq (k-1)!(1 + o_n(1)) \log_2 n$  (Entropy method).
- $N^*(n, k) \leq \frac{k}{\log_2 \frac{k!}{k!-1}} \log_2 n + 1$  (Probabilistic method).
- Fix  $(j_1, \dots, j_k) \in S_{n,k}$ .  $\Pr[\neg(\pi(j_1) < \dots < \pi(j_k)) \mid \pi \stackrel{u}{\leftarrow} S_n] = 1 - \frac{1}{k!}$
- Sample indep. and unif. at random  $\mathcal{F}_d = \{\pi_1, \dots, \pi_d\} \leftarrow S_n$
- $\Pr[\forall \pi \in \mathcal{F}_d : \neg(\pi(j_1) < \dots < \pi(j_k))] = (1 - \frac{1}{k!})^d$
- Hence,

$$\begin{aligned} \Pr \left[ \bigcup_{(j_1, \dots, j_k) \in S_{n,k}} \forall \pi \neg(\pi(j_1) < \dots < \pi(j_k)) \right] &\leq \sum_{(j_1, \dots, j_k) \in S_{n,k}} \left(1 - \frac{1}{k!}\right)^d \\ &= \binom{n}{k} k! \left(1 - \frac{1}{k!}\right)^d \stackrel{!}{<} 1 \end{aligned} \quad (*)$$

- $\Rightarrow$  smallest  $d$  ensuring  $(*)$  guarantees existence of SCA on  $d$  rows

## Definition (Yuster 2020)

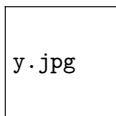
A strength- $k$  SCA  $\mathcal{A} \subseteq S_n$  is called  $\lambda$ -perfect, if each  $x \in S_{n,k}$  is covered  $\lambda$  times in  $\mathcal{A}$ . Notation:  $\mathcal{A} \in \text{PSCA}(n, k, \lambda)$ .

- $\mathcal{A} \in \text{PSCA}(n, k, \lambda) \implies |\mathcal{A}| = \lambda k!$ .
- Set  $g(n, k) := \min\{\lambda : \text{PSCA}(n, k, \lambda) \neq \emptyset\}$ .
- Levenshtein 1992:  $g(n, n-1) = 1$  as  $\mathcal{A} \in \text{PSCA}(n, n-1, 1)$  constructable

## Theorem (Yuster 2020)

Lower bound  $g(n, k) = \Omega(n^{k/2})$  applies.

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Raphael Yuster: Can we find “small”, i.e., polynomial size PSCAs?

- Master's thesis: Yes, always (poly. size:  $g(n, k) \leq \frac{1}{k!} n^{O(k^2/\ln k)}$ ).

# Fascination PSCAs – when optimal packing succeeds...

PSCA(4, 3, 1)	PSCA(5, 3, 1)	PSCA(5, 3, 2)
<pre> 3 1 2 4 1 3 4 2 2 3 4 1 4 1 2 3 2 1 4 3 4 3 2 1                     </pre>	$\emptyset$	<pre> 1 2 4 5 3 1 3 4 5 2 2 1 5 4 3 2 3 4 1 5 3 1 5 4 2 3 2 5 1 4 4 1 3 2 5 4 2 3 5 1 4 5 2 1 3 5 1 2 3 4 5 3 2 4 1 5 4 3 1 2                     </pre>
$\Rightarrow g(4, 3) = 1$	$\Rightarrow g(5, 3) > 1$	$\Rightarrow g(5, 3) = 2$

A construction of Tarui et al. 2000 can be slightly optimized:

$n (k = 4)$	A	B	$n (k = 3)$	A	B
5	1	1	4	1	1
6	1	672	5-7	2	10
7	2	672	8	3	-
8-12	18	-	9	4	-
13	234	-	10-12	6	-
14-17	5040	<b>672</b>	13-14	7	-
18-21	5040	3 139 584	15-16	16	<b>10</b>
22	18 480	-	17-19	19	180
23	425 040	-	20-32	96	180
24	10 200 960	<b>3 139 584</b>	33-64	/	180
25-65	/	3 139 584	65-256	/	340
66-257	/	23 482 368			

Upper bounds for  $g(n, k)$ :

A  $\rightsquigarrow$  State of the art (Na et al. 2022, Gentle et al. 2022);

B  $\rightsquigarrow$  Recursive construction of Tarui et al. 2000, optimized (Iurlano 2022)

Green: optimal bounds

- Mathon et al. 1999:  $\lambda = 1$ : Backtracking within unions of cosets of  $S_n$
- Na et al. 2022:  $\lambda$  arbitrary: analogous backtracking within unions ...

## Example (Na et al. 2022)

Subgroup  $H := \{id, h_2, h_3\} \leq S_7$ ,  $\alpha, \beta, \gamma, \delta \in S_7$ , such that the union of cosets

$$\mathcal{A} := H\alpha \cup H\beta \cup H\gamma \cup H\delta \in \text{PSCA}(7, 3, 2).$$

Four left cosets of the order 3 subgroup  $\langle 3275641 \rangle =: H$

$\langle \sigma \rangle$	1573426 $\langle \sigma \rangle$	4261735 $\langle \sigma \rangle$	4756123 $\langle \sigma \rangle$
1234567	1573426	4261735	4756123
3275641	3617524	5243176	5164327
7216453	7431625	6257314	6345721

*Handwritten notes:*  $H = \langle id, id\sigma, id\sigma^2 \rangle$  with arrows pointing to the first column of the table. A red arrow points from the bottom row of the table to the right.

- Gentle et al. 2022: Analysis of columnwise distribution of symbols
- Master's thesis: Study of feasible gaps between  $\ell$ -subsequences ( $\ell < k$ )
  - $\rightsquigarrow$  Integer Linear Program formulation for "admissible gap configs"
  - $\rightsquigarrow$  Particular solution of ILP "reflection symmetry"



Reflection symmetry illustrated:

1	2	3	4	5	6
1	4	3	2	6	5
2	5	6	1	3	4
3	5	1	6	2	4
3	6	1	5	4	2
4	6	5	1	2	3
6	4	2	3	1	5
6	3	2	4	5	1
5	3	4	2	1	6
5	2	4	3	6	1
4	1	5	6	3	2
2	1	6	5	4	3

1	2	3	4	5	6
1	4	3	2	6	5
2	5	6	1	3	4
3	5	1	6	2	4
3	6	1	5	4	2
4	6	5	1	2	3
6	4	2	3	1	5
6	3	2	4	5	1
5	3	4	2	1	6
5	2	4	3	6	1
4	1	5	6	3	2
2	1	6	5	4	3

- Complexity class of calculation of  $N^*(n, k)$  resp.  $g(n, k)$ ?
- Correctness of the Levenshtein conjecture

$$k \notin \{1, 2, 4\} \Rightarrow g(k + 2, k) > 1?$$

- Does even  $\text{PSCA}(n, k, \lambda) \neq \emptyset \Rightarrow \text{PSCA}(n, k, \lambda + 1) \neq \emptyset$  apply?
- Calculate next value  $g(9, 3) \in \{3, 4\}$ .
- Similar combinatorial structures: Does union-of-cosets-approach succeed?

## Shortest Superpermutation Problem

- Motivation: Watch  $n$  episodes of a series in all  $n!$  possible ways (contiguously) by a long chain of movie events in order to have “viewed it correctly at least once”
- $n \in \mathbb{N}$ . Find shortest length- $L$  string  $S$  such that

$$\forall x \in S_{n,n} : x \text{ substring } S.$$

- e.g.  $n = 4$ : 123412314231243121342132413214321 ( $L = 33$ , 1432 contained)
- Solving approach: [open dedicated pdf](#)

- Problem: Find shortest *Universal Cycle Covering for given  $k$ -Subset* (details see Dębski et al. 2016)
- A cyclic sequence of elements of  $[n]$  is an  $(n, k)$ -*Ucycle covering* if every  $k$ -subset of  $[n]$  appears in this sequence at at least once as a subsequence of consecutive terms.
- Problem more attackable ... not “doubly exponential” effort (cf. Superpermutation problem)
- Ideas from <https://www.ac.tuwien.ac.at/files/pub/mayerhofer-22.pdf>

### A Beam Search for the **Shortest Common Supersequence Problem** Guided by an Approximate Expected Length Calculation










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exploitable?

Nice to find:

- Idea good heuristic/construction for (near-/ optimal) PSCAs / SCAs
- Reduce an NP-hard problem to  $\text{CALC}(N^*(n, k))$  or  $\text{CALC}(g(n, k))$

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